

Resiliency in Ethernet Based Transport Networks

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Outline

- Introduction
- What is “switched” Ethernet?
- Legacy Ethernet
- Security and Reliability issues
- Rapid spanning tree protocol & others
- Multiple spanning trees
- Carrier-grade Ethernet
- Ethernet layer 2 switching
- Conclusions

Introduction

- The topic of this presentation is resiliency in native Ethernet transport networks.
 - Ethernet services (e.g., E-Line, E-LAN) can be based on virtually any transport technology like EoS, IP/MPLS VPNs, ATM...
 - Resiliency in Ethernet services depends on underlying transport network
 - In p-to-p links between IP routers or such, Ethernet is just one LL protocol
 - “Carrier-grade Ethernet” sometimes mean Ethernet over MPLS or such
- Ethernet is considered to be cheap, simple and versatile transport technology for IP NGN
- Ethernet is also considered to be easy to manage, plug-and-play solution
- Most of the Internet hosts use Ethernet i/f — so why not transport that traffic end-to-end in native form?

What is “switched” Ethernet?

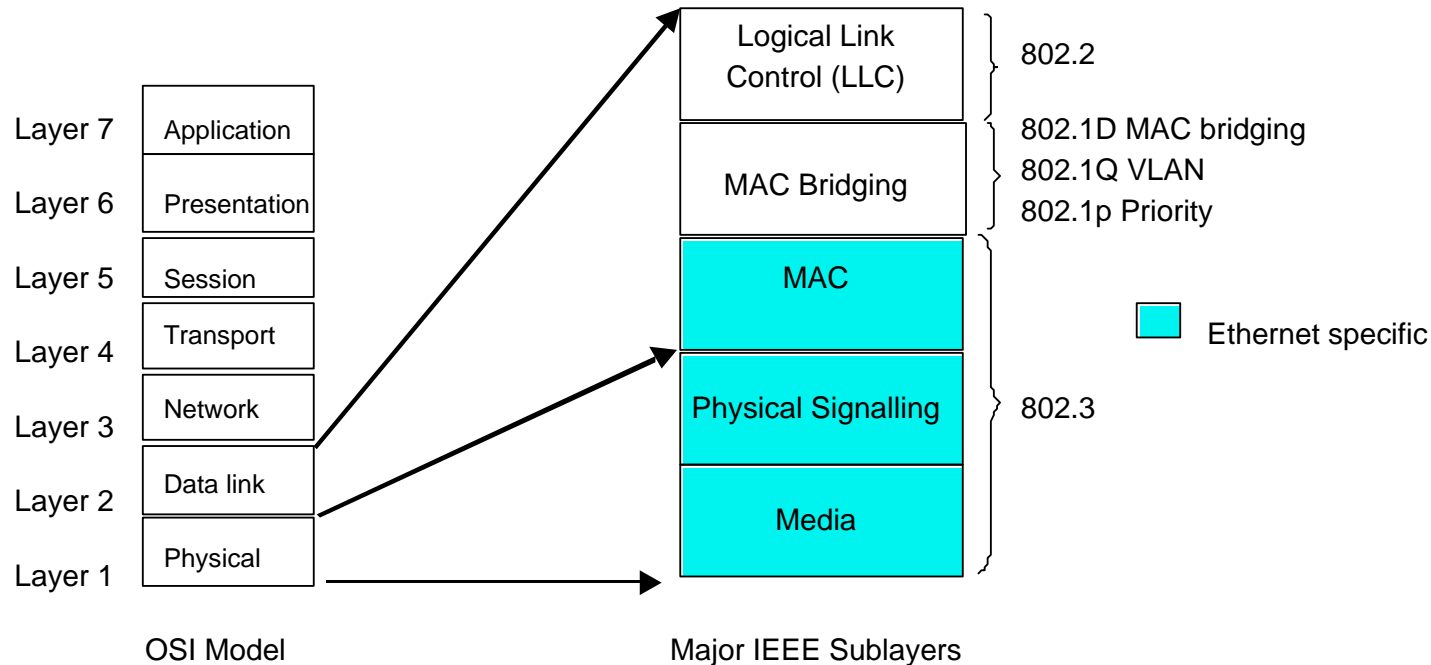
Definition of layer 2 protocol:

Anything defined by a committee whose charter is to standardize a layer 2 protocol. (*Radia Perlman, Sun Laboratories*)

- Ether**NET** is a network — right?
 - No — it’s just a link in a network.
- But there are Ethernet switches that do L2 switching?
 - Not exactly — logically an Ethernet switch is more or less a set of bridges connected by a high-speed Ethernet segment (or multiport bridge).
 - Bridged/switched Ethernet divides a single collision domain (Eth. segment) to multiple ones (rather than divides a network into multiple subnets)

Legacy Ethernet

Ethernet is a frame based multiaccess LL protocol with flat addressing.



- Bridge learning make it possible to divide a LAN into multiple collision domains automatically
- Need loop-free topology → Spanning Tree Protocol (STP)

Security, Reliability and Scalability issues

- Performance monitoring & fault detection
 - Only 10GbE WAN option (poor man's SDH) has any performance monitoring
 - Otherwise — only “loss-of-light” at Ethernet layer
 - How to get performance data for maintenance, e.g., replacing aging lasers?
- Management & monitoring
 - Relies on upper layer (IP) system like SNMP or Web application
- STP
 - Root node is a single point of failure
 - Sloooooow action in error situations (i.e., up to minutes)
 - Insecure — if user's port accepts STP packets (e.g., due to misconfig.), attacker can bring down the network or make all traffic go through attacker's machine

- No load balancing (except for link bundles)
- Link aggregation — parallel links between two switches can be bundled to a single virtual link for STP
 - Provides link-level load sharing
 - Std. version provides “fast failover” — ≈ 1 s
 - faster recovery with vendor specific solutions
- Cheap shall it be...
 - 10 GbE WAN option is much like SDH/SONET but eliminates “expensive” SDH clock distribution...
 - ...actually money saved was about 5 USD per line card
 - 10 GbE transport network cannot be used to distribute clock signal to legacy networks

Rapid spanning tree protocol & others

- STP relies on timers and Bridge Protocol Data Units (BPDU) sent by root node — in RSTP (802.1w) each bridge sends hello BPDUs → faster fault detection
- All edge ports can be turned on immediately
- If a non-root bridge detects that root is unreachable, it can immediately declare itself as a new root node
- Still ≈ 1 s convergence
- Rapid Ring STP (RRSTP) or Ethernet Automatic Protection Switching (EAPS) for bidirectional Ethernet rings (\neq RPR) — mimics bidir. SDH ring with similar performance

Multiple spanning trees

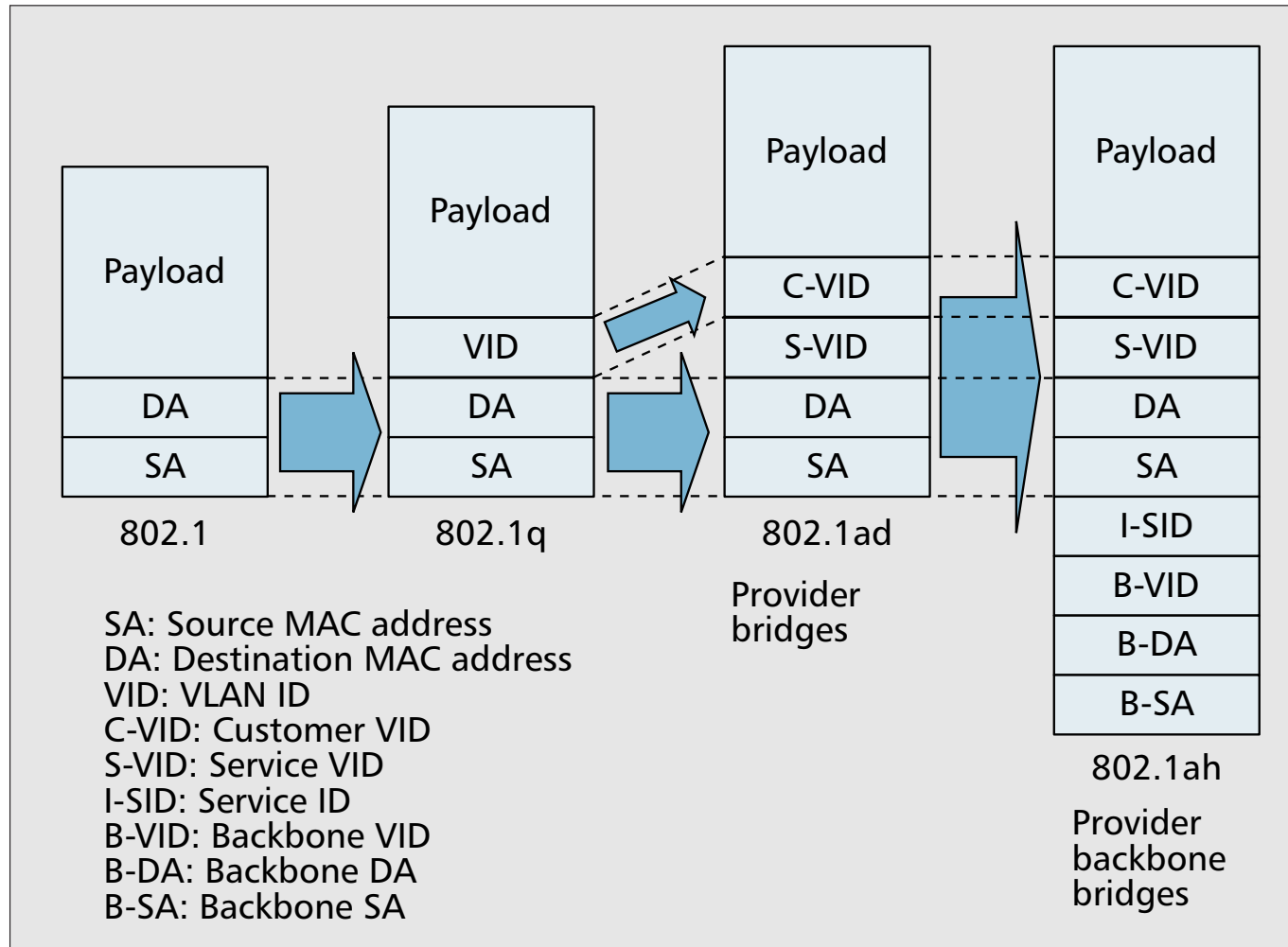
- RSTP has still only one spanning tree
 - Suboptimal “routes” for most of VLANs
 - Bottleneck links and unused resources
- MSTP (802.1s) allows for own spanning tree for each VLAN
 - Location of root node for each ST can be selected according to each structure/traffic patterns of each VLAN → better “routes”
 - Traffic is more distributed and resources are better used
 - A single fault does not (necessarily) affect all VLANs

Carrier-grade Ethernet

Major requirement: the transport network has to be protected from user's actions
→ hierarchical network is needed

- Service OAM (802.1ag connectivity fault management)
 - Continuity check, link trace, loopback and AIS messages
- Link OAM (802.3 clause 57)
 - Discovery, link monitoring, RFI, remote loopback
- L2 VPN is considered as a major application
 - Q-in-Q does not provide sufficient protection for transport network, MAC learning is also an issue
 - Limited number (4094) of VLANs
 - MAC-in-MAC is required (that is, user's frames are carried inside provider's frames)

- TDM over Ethernet to support legacy PDH services



Ethernet L2 switching

(a.k.a. Independent VLAN learning (IVL))

- Even with MSTP exact traffic engineering (“route pinning”) is not possible
- L2 switching with Ethernet
 - VLAN ID (VID) and destination MAC addr. are used as a label (VID can be swapped, MAC addr. has to remain the same)
 - Ethernet LSP (or PVCCs) are created by configuring label lookup tables in switches
 - Only part of VID range can be used
- Operation: if VID is in the range of Ethernet switching
 - output port, new VID = Lookup(VID, MAC addr.)

Conclusions

So, is Ethernet up to carrier-grade transport network standards?

- With all the modifications: probably it is, but I doubt that it would be simple, cheap and easy-to-manage anymore...
- Fault-tolerance
 - Restoration: with RSTP & MSTP 1 s timescale and no single point of failure — comparable to PNNI
 - Link and ring protection: with tuned-up link bundles and RRSTP comparable to SDH
 - Path protection: ?
- “Cell-tax” was big issue — but does anyone care about efficiency anymore?
 - Voice over IPv6 over MPLS over MAC-in-MAC Ethernet (i.e. over 100 octets overhead per packet...)?

RFC 1925 (The Twelve Networking Truths)

...

- (3) With sufficient thrust, pigs fly just fine. However, this is not necessarily a good idea. It is hard to be sure where they are going to land, and it could be dangerous sitting under them as they fly overhead.

...

- (6) It is easier to move a problem around (for example, by moving the problem to a different part of the overall network architecture) than it is to solve it.

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- (11) Every old idea will be proposed again with a different name and a different presentation, regardless of whether it works.

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